

Optimal parameter-space covering for continuous-wave searches

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GWDAAW11@Potsdam Dec 19, 2006

Matched filtering: parameter-space covering

Templates parametrized by $\lambda \equiv \{\lambda_1, \lambda_2, \dots, \lambda_n\} \in \mathbb{P}_n$

Detection statistic: $\mathcal{F}(\lambda)$.

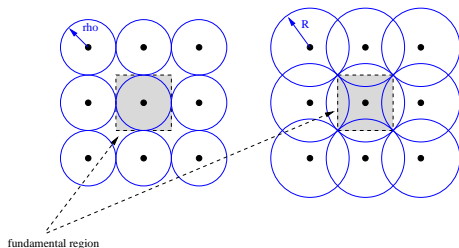
“Distance” $m \equiv \frac{E[\mathcal{F}(\lambda_s) - \mathcal{F}(\lambda)]}{E[\mathcal{F}(\lambda_s)]} = g_{ij}(\lambda_s) \Delta \lambda^i \Delta \lambda^j + \mathcal{O}(\Delta \lambda^3)$

Optimal template bank:

- Set of “templates” $\lambda_{(i)}$ such that for ANY point $\lambda \in \mathbb{P}_n$, the mismatch m to the *nearest* template satisfies $m \leq m_{\max}$.
 \implies spheres of “covering radius” $R = \sqrt{m_{\max}}$.
- Find the covering $\{\lambda_{(i)}\}$ with the **smallest** number of templates \implies “sphere covering problem” (‘dual’ of sphere-packing problem!)

Sphere “Packing” versus “Covering”

Conway, Sloane, *Sphere packings, lattices and groups* (1998)



Packing density: $\Delta \equiv \frac{\text{Volume of sphere}}{\text{fundamental volume}} < 1$

☞ fraction of space occupied by spheres

Covering thickness: $\Theta \equiv \frac{\text{Volume of sphere}}{\text{fundamental volume}} > 1$

☞ average number of spheres covering a point

Packing problem: given ρ , maximize density Δ

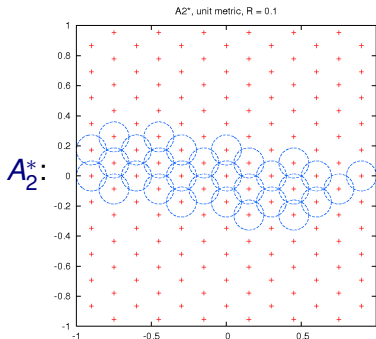
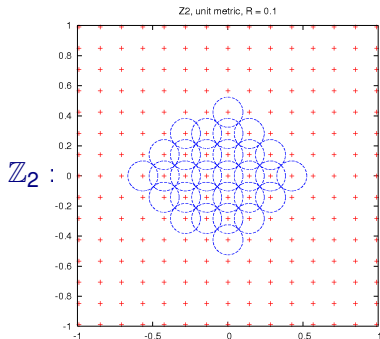
Covering problem: given R , minimize thickness Θ

What is know about (Euklidean) Covering?

n	lattice	Θ	alt.	Θ
2	A_2^* (hexagonal)	1.21		
3	A_3^* (bcc)	1.46		
4	A_4^*	1.77		
5	A_5^*	2.12		
6	L_6^C	2.46	A_6^*	2.55
7	L_7^C	2.90	A_7^*	3.06
8	L_8^C	3.14	A_8^*	3.67
9	L_9^C	4.27	A_9^*	4.39

- Today L_n^C best covering known for up to $n = 15$
- Generally A_n^* (previous record-holder) not much thicker than the best currently known lattice.

Advantage of A_n^* covering over \mathbb{Z}^n



“gain” of A_n^* with respect to \mathbb{Z}_n : $\kappa(n) = \frac{\Theta(\mathbb{Z}_n)}{\Theta(A_n^*)} \sim e^n$

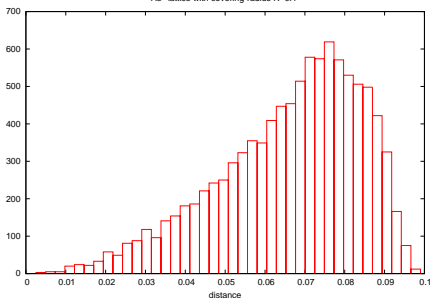
n	2	3	4	5	11	17
$\kappa(n)$	1.30	1.86	2.80	4.33	78.23	1691.6

Monte Carlo - check of covering-radius

A_n^* lattice covering implemented in `LALLatticeCovering()`

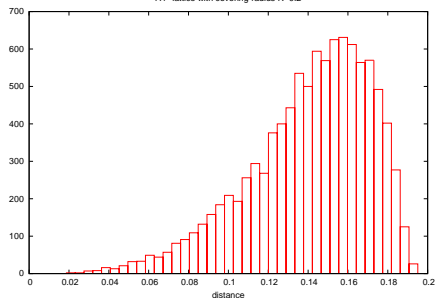
- generate A_n^* lattice for given covering radius $R = \sqrt{m_{\max}}$
- pick random points and determine their minimal match m
- 👉 histogram m

A3*-lattice with covering-radius R=0.1



$n = 3, R = 0.1$

A4*-lattice with covering-radius R=0.2



$n = 4, R = 0.2$

Flat metric approximation

We need **explicitly** flat metric to use lattice covering, but the continuous-wave metric g_{ij} in $\lambda^i = (\alpha, \delta, f, \dot{f}, \dots)$ is **curved**.

- Approx. I: “phase metric”: $g_{ij} \sim g_{ij}^\phi = \langle \partial_i \phi \partial_j \phi \rangle - \langle \partial_i \phi \rangle \langle \partial_j \phi \rangle$

$$\phi(t) = 2\pi \left[f \left(t + \frac{\mathbf{r}(t) \cdot \mathbf{n}}{c} \right) + \frac{1}{2} \dot{f} \left(t + \frac{\mathbf{r}(t) \cdot \mathbf{n}}{c} \right)^2 + \dots \right]$$

- Approx. II: “JKS simplified phase”: (linear Roemer delay)

$$\phi(t) \approx 2\pi \left[f t + \frac{1}{2} \dot{f} t^2 + \dots + \frac{\mathbf{r}(t) \cdot \mathbf{n}}{c} f(t) \right]$$

- Approx. III: “Orbital metric”: neglect $\mathbf{r}_{\text{spin}}(t)$: (true for LISA)

$$\phi_{\text{orb}}(t) \approx 2\pi \left[f t + \frac{1}{2} \dot{f} t^2 + \dots + \frac{\mathbf{r}_{\text{orb}}(t) \cdot \mathbf{n}}{c} f(t) \right]$$

☞ g_{ij}^{orb} can be shown to be **flat!**

Plans

- Improve LatticeCovering() to generate lattice “point-by-point” instead of “all-at-once”
- Use flat-metric approximation to build A_n^* covering for continuous-wave searches (LIGO and LISA)
- Test this by Monte-Carlo injection and signal recovery
- Alternative approach [with X. Siemens]: use (strong) degeneracy of orbital metric to reduce the “effective dimensionality” of the parameter space